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DOI [10.1109/TITS.2019.2914770](https://doi.org/10.1109/TITS.2019.2914770)

Publication date 2019 Document Version

Final published version

Published in IEEE Transactions on Intelligent Transportation Systems

Citation (APA)

Sun, J., Vû, H., Ellerbroek, J., & Hoekstra, J. (2019). pyModeS: Decoding Mode-S Surveillance Data for Open Air Transportation Research. IEEE Transactions on Intelligent Transportation Systems, 21(7), 2777- 2786. Article 8718517. <https://doi.org/10.1109/TITS.2019.2914770>

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pyModeS: Decoding Mode-S Surveillance Data for Open Air Transportation Research

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*Abstract***— The availability of low-cost automatic dependent surveillance-broadcast (ADS-B) receivers has given researchers the ability to make use of large amounts of aircraft state data. This data is being used to support air transportation research in performance study, trajectory prediction, procedure analysis, and airspace design. However, aircraft states contained in ADS-B messages are limited. More performance parameters are downlinked as Mode-S Comm-B replies, upon the automatic and periodic interrogation of air traffic control secondary surveillance radar. These replies reveal aircraft airspeed, turn rate, target altitude, and so on. They can be intercepted using the same 1090-MHz receiver that receives the ADS-B messages. However, a third-party observer does not know the interrogations, which originated the Comm-B replies. Thus, it is difficult to decode these messages without knowing the type and source aircraft. Furthermore, the parity check also cannot be performed without knowing the interrogations. In this paper, we propose a new heuristic-probabilistic method to decode the Comm-B replies and to check the correctness of the messages. Based on a reference dataset provided by air traffic control of the Netherlands, the method yields a success rate of 97.68% with an error below 0.01%. The performance of the proposed method is further examined with data from eight different regions of the world. The implementation of the inference and decoding process,** *pyModeS***, is shared as an open-source library.**

*Index Terms***— Aircraft surveillance, air traffic control, Mode-S, ADS-B, Comm-B, enhanced Mode-S.**

I. INTRODUCTION

I^N air transportation research, studies related to aircraft
performance are often dependent on the airspeed of the
signals. This good information is used in the dynamic model N air transportation research, studies related to aircraft aircraft. This speed information is used in the dynamic model of the aircraft to perform, for example, state estimations [1] and trajectory predictions [2]. In addition to airspeed, the performance model also takes into account other trajectory state information, such as positions, ground speeds, and altitudes.

Many of these states in the dynamic model can be openly observed using modern aircraft surveillance technology, for instance, the *Automatic Dependent Surveillance-Broadcast* (ADS-B). ADS-B provides information on aircraft position, speed, and vertical rate. The speed contained in ADS-B refers to the ground speed rather than the airspeed, unless in (rare) cases when the location cannot be determined from the Global

Digital Object Identifier 10.1109/TITS.2019.2914770

Navigation Satellite Systems. The advantage of ADS-B is that the signals can be openly intercepted and decoded using a simple ground receiver set-up.

When airspeed is not available, there are two ways to adopt the ground speed for performance analysis. The first simple approach is to assume the ground speed as airspeed by ignoring the wind. This may cause errors in performance calculations when a strong wind is present. The second approach is to integrate the wind from numerical forecast models. However, wind data from these models often cannot reflect local wind variations accurately.

On the other hand, air traffic controllers are also interested in the same performance parameters. These parameters are constantly interrogated by surveillance radars under the *Mode-S Enhanced Surveillance* (EHS) technology. Corresponding messages are downlinked using the *Comm-B* protocol. Air traffic controllers make use of this data to better monitor and predict flights and to make better traffic control decisions. Within these downlinked messages, information such as true airspeed, indicated airspeed, Mach number, and true heading of the aircraft are transmitted.

If available, direct access to air traffic control data would provide the most accurate information. However, due to licenses and data agreement processes, obtaining this data can be challenging for third-party researchers. Even when the access is granted, the information is often extracted from historical data archives, which makes it difficult to perform real-time performance analysis.

Nevertheless, it is possible to obtain the downlinked Comm-B data openly just as ADS-B data with the same ground receiver. However, many difficulties arise when one tries to decode these reply messages. The biggest barriers for decoding are the unknown aircraft source represented by the ICAO transponder address and the interrogation type represented by the *Comm-B Data Selector* (BDS) code. Even though the structure of messages follows open standards [3], [4], without knowing the ICAO code and BDS type, one cannot extract useful information from these messages.

The goal of this paper is to enable open and real-time access to these Mode-S messages. The main research questions of this study are defined as:

- 1) How to determine the source aircraft of a Comm-B message?
- 2) How to identify the BDS type and decode a Comm-B message without knowing the original interrogation?

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Manuscript received September 21, 2018; revised March 1, 2019; accepted April 30, 2019. Date of publication May 20, 2019; date of current version June 29, 2020. The Associate Editor for this paper was Z.-H. Mao. *(Corresponding author: Junzi Sun.)*

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Fig. 1. The Mode-S inference and decoding pipeline.

3) How to detect errors in Comm-B messages under incomplete information?

To illustrate these inference efforts, Fig. 1 shows the process related to the methods that are proposed in this paper. In this figure, the downlink surveillance signal (containing ADS-B and Comm-B) from aircraft is received by a software-defined radio (SDR) receiver first. The signal is converted to a raw binary data stream, which is then further decomposed into a sequence of message frames. For ADS-B messages, information can be decoded directly. For Comm-B replies, we first use the inference methods proposed in this paper to determine the BDS code, source aircraft, and errors. Finally, the information contained in these messages is decoded.

In the remaining sections of this paper, we first explain the background related to ADS-B and EHS. Next, we discuss the identification processes and error detection in detail. Several tests are proposed in this paper. The methods are also validated with a reference dataset provided by Air Traffic Control the Netherlands (LVNL). Finally, we discuss the use cases, implementation, and recommendations, as well as the conclusions.

In addition to the identification and decoding process described in this paper, an open-source decoder library, *pyModeS*, which is implemented using Python programming language has also been made public.¹

II. BACKGROUND

A. Fundamentals of Mode-S, ADS-B, and Comm-B

As the demand for air transportation increases, airspace over the world is becoming more crowded. To efficiently make use of the airspace and increase the traffic capacity, air traffic controllers need to rely on accurate flight trajectory predictions. Communications between ground and aircraft are becoming more frequent in order to obtain accurate updates regarding aircraft states beyond the basic position that is provided by the primary radar. Since the beginning of aircraft surveillance, different methods and standards have been developed for downlinking aircraft data. The two most common methods are based on interrogation and broadcast.

Traditionally, when an air traffic controller requires information in addition to the aircraft position, the Mode-S selective interrogation [5] is used. This is performed by the *secondary surveillance radar* (SSR). Numerous aircraft states can be interrogated by the SSR. The most common downlinked messages are Comm-B replies. The content of interrogation is identified by the BDS code, which is a twodigit hexadecimal code (8 bits) that indicates the information desired by the air traffic controller. In total, 255 BDS codes can be defined. The reply data is encoded in a 112-bit Comm-B downlink message. Among all these BDS codes, several BDS codes are grouped and identified as Mode-S EHS, which consists of *selected intention report* (BDS 40), *track and turn report* (BDS 50), and *heading and speed report* (BDS 60).

The simpler ADS-B is an implementation of Mode-S extended squitter [4]. It is a newer technology compared to the interrogation-based Mode-S, which allows the automatic broadcast of the aircraft state information at a constant rate. In many regions, aircraft are required to be equipped with Mode-S transponders that are compatible with ADS-B. When it is enabled, ADS-B allows aircraft to automatically report the identification, location, speed, and operational status. The update interval of critical states (such as position and speed) is designed to be around 0.5 seconds.

Both ADS-B messages and Comm-B replies are transmitted using the 1090 MHz transponder. Downlinked signals can be intercepted freely using low-cost commercial off-the-shelf ground receivers. Several crowd-sourced initiatives have been constructing global networks of ground receivers, for example, *ADS-B Exchange*, *FlightRadar24*, *FlightAware*, and *OpenSkynetwork*. The quantity of data gathered by these networks is enormous, which leads to great potential in air transportation researches. For example, in recent research, this data has been used for operational performance studies [6] and trajectory prediction [7]. The ground receiver networks also enable the possibility to determine aircraft location by using multilateration [8].

ADS-B is designed as an independent communication protocol, where the message itself contains all information needed for decoding. On the contrary, Comm-B communication is designed as a dependent protocol. Only the air traffic controller who initiated the interrogation can identify the source aircraft and decode the content of the replies. To this extent, third-party observers have no information on the interrogated aircraft or corresponding BDS code. However, it is has been shown that some information is possible to be extracted in an earlier research [9], which is used to provide meteorological observations.

¹Available at: https://github.com/junzis/pyModeS

Fig. 2. The structures of ADS-B and Mode-S Comm-B messages.

B. Regulation and Availability

Several Mode-S capabilities are mandatory for aircraft flying in European airspace since 2009. Two different categories of Mode-S surveillance are defined, which are elementary surveillance (ELS) and enhanced surveillance (EHS) [10]. According to European regulation, all aircraft that fly Instrument Flight Rules (IFR) in general air traffic (GAT) must be ELS compliant. In addition, all fixed-wing aircraft flying IFR in GAT with a maximum take-off mass greater than 5.7 ton or a maximum cruising true airspeed greater than 250 knots must be EHS compliant. ADS-B is a newer surveillance technology and has also been adopted broadly. Regulators in both Europe and the United States have set the agenda for obligatory compliance.

Since ADS-B does not require active interrogations from surveillance radar, the messages are broadcast and available at all times everywhere. These messages can be received by ground receivers or satellites [11].

In Mode-S ELS, only a limited number of parameters are reported, including aircraft identity, altitude, flight status, and related supporting parameters. In Mode-S EHS, more aircraft states are interrogated, such as indicated airspeed, Mach number, vertical rate, magnetic heading, track angle, roll angle, selected altitude, and ground speed.

Depending on the location of the (third-party) ground receiver, the number of received replies varies. The availability and quantity of Comm-B messages also depend on air traffic density and the number of secondary surveillance radars in the area, as well as the rate of interrogation.

C. Data Structure

In this paper, we focus on two types of messages, which are ADS-B message and Comm-B message from Mode-S EHS. The structures of ADS-B and Comm-B messages are defined in ICAO Annex 10 [12]. ADS-B and Mode-S data are constructed using the data frame shown in Fig. 2, with a total message length of 112 bits. The number of bits of each segment is indicated with parentheses in this figure. Each message starts with the *downlink format* (DF), followed by a 27-bit header with different components. Then, the crucial 56-bit data is appended with the downlink information encoded. Lastly, 24 bits are dedicated to the parity checksum.

ADS-B messages are identified by a DF number of 17 (10001 in binary format). In the header of an ADS-B message, the address of the aircraft transponder is indicated. This is a 24-bit address assigned by ICAO and categorized according to geographic region and country. The leading 3 bits are sub-type or category in different types of ADS-B messages.

Fig. 3. Number of ADS-B message and Mode-S Comm-B replies received in 24 hours, May 30, 2018.

The *Type Code* (TC) is set using the first 5 bits of the 56-bit data segment. It defines the general type of message, for example, airborne position, airborne velocity, surface position, identification, etc.

In a Comm-B reply message, the DF number can be either 20 or 21 (10100 or 10101 in binary format). In the case of DF = 20, the last 15 bits of the header indicates the *Altitude Code* (AC). When $DF = 21$, the last 15 bits represent the *Identification Code* (ID) (a.k.a: the *squawk code*). The leading three segments in the header are *Flight Status* (FS), *Downlink Request* (DR), and *Utility Message* (UM). Unlike ADS-B, there is no indication of ICAO transponder address nor the BDS code in a message, except for a few cases.²

Fig. 3 illustrates the distribution of ADS-B and Comm-B replies, together with their distinct types, for a 24-hour period of data from a ground receiver situated in Delft, the Netherlands. We can see there are more Comm-B messages than ADS-B messages. About 35 out of 38 million EHS messages are not directly identifiable. Unlike ADS-B, none of the Comm-B messages can be checked for corruption due to the incomplete information on aircraft source and BDS code.

It is worth pointing out that more than half of the ADS-B messages (16 out of 28 million messages) are corrupted. The corruption of messages is also analyzed in the later sections of this paper.

III. BASIC DECODING OF ADS-B DATA

As mentioned earlier, the content of an ADS-B message is directly identified by its Type Code. The primary parameters downlinked by each type are listed in Table I.

Within the same group, Type Code values also have their own indications. For example, in the identification group, a TC of 1 to 4 indicates different aircraft categories defined by the Mode-S standard according to [3]. In position groups, TC defines the level of accuracy and uncertainty.

Since ADS-B was first introduced up until now, there have been different versions of implementation: version 0,

²The exception cases are BDS 10, 20, and 30 messages. The BDS code is indicated in these messages. However, together, they only represent a very small percentage of all messages.

TABLE I BASIC ADS-B MESSAGE TYPES (VERSION 0, 1, 2)

TC	Content	Primary parameters				
$1 - 4$	Aircraft identification	Call sign				
$5 - 8$	Surface position	Latitude				
		Longitude				
		Speed				
		Track angle				
$9 - 18$	Airborne position	Latitude				
$20 - 22$		Longitude				
		Altitude				
19	Airborne velocity	East-west component				
		North-south component				
		Vertical rate				

TABLE II ADDITIONAL ADS-B MESSAGE TYPES (FOR VERSION 1, 2)

version 1, and version 2. In version 1 and 2, the accuracy and uncertainty are re-classified using two additional supplement bits. They offer more levels of *Navigation integrity category* (NIC). In addition, from ADS-B version 1 onward, three more types are introduced, as shown in Table II.

The version number 1 or 2 can be retrieved in $TC = 31$ messages. However, for version 0 transponder, no information regarding the version number is transmitted. In order to identify the exact version of the ADS-B transponder, we have to first check whether the aircraft is broadcasting any message with TC of 28, 29 or 31. If not, then the ADS-B version is assumed to be version 0. When a TC 31 message is received, the ADS-B version (1 or 2) can be found at the $73th$ to $75th$ bit of the message.

Knowing the Type Code, in combination with the ADS-B version, it is possible to identify the uncertainties in broadcast aircraft states. At the same time, payload information can be decoded according to the *Technical Provisions for Mode Service and Extended Squitter*, which is published by ICAO [3], [4]. Some of the ADS-B states, such as ICAO address and speed, are used for the Comm-B BDS inference in this paper.

IV. INFERRING THE COMM-B DOWNLINK PARAMETERS

Without knowledge of the original interrogation request, most of the Comm-B replies cannot be identified directly due to the lack of BDS information. It is also not possible to know

TABLE III COMMON MODE-S COMM-B BDS CODE AND CONTENT

Type	BDS	Content	Parameters			
ELS	10	Data link capability	Transponder capabilities			
	17	GICB capability	Mode-S data capabilities			
	20	Aircraft identification	Call sign			
	30	ACAS resolution	Resolution advisories			
			Threat identity			
EHS	40	Vertical intention	Select altitude			
			Barometric pressure setting			
	50	Track and turn	Roll angle			
			Track angle			
			Ground speed			
			Track angle rate			
			True airspeed			
	60	Heading and speed	Magnetic heading			
			Indicated airspeed			
			Mach number			
			Vertical velocity			

whether a message is corrupted because the source aircraft ID (ICAO address) is unknown. This is because the checksum in Comm-B reply is overlaid with the 24-bit ICAO address. In some cases, it is overlaid again with the BDS code.

The two-digit hexadecimal BDS code can support up to 255 different types of messages. In practice, only a small portion of these types are interrogated. According to the European ELS and EHS mandate, the most commonly interrogated BDS codes are 10, 17, 20, 30, 40, 50, and 60. In Table III, parameters contained in these reply messages are listed.

In order to decode these messages, we first need to recover the ICAO address of the source aircraft and examine whether the messages are corrupted without full knowledge of parity. In parallel, we also implement the process of identifying the BDS codes, which are used for further error detection. As a result, the decoding of a message can be accomplished.

A. Source Aircraft Identification and Error Detection

Thanks to a process called *Address Parity* [13], we can recover the ICAO addresses for most of the Comm-B messages. In the last 24 bits of the message, a parity checksum is inserted. The checksum is computed using a *cyclic redundancy check* (CRC) algorithm [14], which is a common errordetecting scheme in telecommunications.

The generator code used for Mode-S CRC encoding is specifically designed in the following binary format:

$$
G = 1111111111111101000001001 \tag{1}
$$

which is used to compute the checksum by the encoder and to validate the message by the decoder. In the polynomial form, the generator is expressed as:

$$
G(x) = x24 + x23 + x22 + x21 + x20 + x19 + x18 + x17 + x16 + x15 + x14 + x13 + x12 + x10 + x3 + 1
$$
 (2)

Similarly, the binary format of the first 88 bits of the message can also be written in the polynomial format with the

Fig. 4. The ICAO address recovery logic.

highest order of 87 (x^{87}) . The message polynomial is denoted as $M(x)$. The checksum is computed using polynomial division between the message and generator. By combining the checksum with parity $P(x)$, we can compute the (possible) ICAO address:

$$
M(x) = \sum_{i=0}^{87} a_i x^i, \quad a_i \in (0, 1)
$$

\n
$$
R(x) = M(x) \% G(x)
$$

\n
$$
A(x) = R(x) + P(x)
$$
\n(3)

where $R(x)$ is the remainder (checksum) of the division $M(x)$ by $G(x)$. In Fig. 4, we illustrate this ICAO address reversal process with an example message.

After the reversal process, the resulting 24 bits would be one of the following possibilities:

- Candidate ICAO address (CA): The correct ICAO address of an aircraft.
- Modified ICAO address (MA): The ICAO address with the first 8 bits overlaid with the BDS code.³ The parity under this condition is identified as *Data Parity*.
- Impossible ICAO address (IA): An ICAO address that is not assigned. This indicates an error in the message.

If corruption occurs in a reply, the reversal process will generate an incorrect ICAO address. This address can still be either CA, MA, or IA. Since we do not know the actual aircraft that has been interrogated, it is not possible to detect the error just using the parity. Instead, we have to identify new methods to identify corrupted messages.

Our proposed error detection mechanism consists of three checks. The first check is to identify impossible ICAO addresses. The unassigned blocks of addresses are listed in Table IV. If the resulting ICAO belongs to one of these blocks, it is likely that the message is corrupted. However, the possibility of Data Parity cannot be ruled out in this case.

Next, the address is cross-validated with the ICAO addresses included in all ADS-B messages from the same time period. Since ADS-B messages can be properly error checked with the CRC process, we are confident about the obtained addresses. After this step, the correct messages are identified. An unidentified ICAO address indicates either corruption in a message or that the aircraft is not equipped with ADS-B capability.

The third error check is the ICAO-to-Squawk comparison. For Comm-B reply message with $DF = 21$ (identity reply),

TABLE IV UN-ASSIGNED ICAO ADDRESS BLOCKS

Starting bits	Address block	Geographic region
00100	200000 - 27FFFF	Africa-Indian Ocean
00101	280000 - 28FFFF	South America
0101	500000 - 5FFFFF	Europe and North Atlantic
01100	600000 - 67FFFF	Middle East
01101	680000 - 6F0000	Asia
1001	900000 - 9FFFFF	North America and Pacific
111011	B00000 - BFFFFF	Caribbean
1101	D00000 - DFFFFF	Reserved
1111	F00000 - FFFFFF	Reserved

Fig. 5. EHS message error identification process.

the squawk code of the transponder can be obtained. By comparing the squawk code and the ICAO address, we can identify error messages, such as the ones associated with low ICAO-to-Squawk combination frequency. Since it is still possible that the message is overlaid with a BDS code, we also use the inferred BDS code to compare the overlaid ICAO address and squawk code.

Summarizing all previous steps, we can construct an error detection model, as illustrated in Fig. 5.

In this model, we first compute three binary scores based on the following conditions:

- 1) The ICAO address is assigned.
- 2) The ICAO address appears in the pool of ADS-B addresses.
- 3) The ICAO address has the ICAO-to-Squawk frequency for more than a certain number of times per minute (six).

The resulting scores are denoted as *s*1, *s*2, and *s*³ respectively. Then the inferred BDS code⁴ is overlaid with the

³Only a small percentage of current aircraft to date are equipped with this capability.

⁴The BDS code is inferred in parallel using the methods described in the following section. This is possible because the deterministic BDS identification can be performed independently.

Fig. 6. Error identification statistics, based on Comm-B replies received at TU Delft from 12:00 to 13:00 UTC, September 7, 2017.

ICAO address to examine the possibility of a Modified ICAO Address. If the MA is applied, the resulting ICAO address with over-lapping BDS code (ICAO/OV) satisfies some of these conditions. Similarly, three more scores are computed, denoted as s'_1 , s'_2 , and s'_3 respectively. The total correctness score of *S* is computed as follows:

$$
S = [(s_1 \vee s_1') \vee (s_2 \vee s_2')] \wedge (s_3 \vee s_3')
$$
 (4)

where ∨ is the logic *OR* operation and ∧ is the logic *AND* operation. The error messages are identified with $S = 0$, while correct messages are identified with $S = 1$. It is important to point out that the error model takes into consideration the situations of 1) aircraft not being equipped with an ADS-B transponder and 2) the Modified ICAO Address is used in a parity checksum.

Fig. 6 shows the resulting decoding statistics based on applying the error detection model to the same one-hour dataset as before. In total, 60% of the messages are identified as correct messages. The remaining 40% are corrupted. We can also see that there are only a few messages that satisfy the condition $s'_2 = 1$, which implies that the BDS overlay is not frequently requested by ATC in our airspace. In addition, only aircraft with overlay capability are able to support this feature. Such information can be found in the BDS 17 (GICB capability) messages of the aircraft.

B. BDS Inference - the Heuristic-Probabilistic Process

Not all aircraft have the same BDS capabilities enabled. To determine which BDS capabilities are available for an aircraft, the SSR would first initiate a Common usage GCIB capability interrogation (BDS 17). There are 24 common BDS capabilities that are reported in the BDS 17 message. SSR will only interrogate the ones that are enabled for this aircraft.

To infer the BDS code request by ATC, a two-step inference process is designed. It consists of a heuristic step and a probabilistic step. The heuristic logic is inspired and developed upon the method proposed by [9], which gives the first estimation of possible BDS codes. The probabilistic identification is introduced to identify messages with both BDS 50 and 60 codes from the previous step. Together, the Heuristic-Probabilistic (HP) process is able to deal with all common ELS and EHS replies.

1) Heuristic Logic: All Comm-B payloads consist of multiple data blocks. Many of these blocks are combined with respective status bits. A status bit must be set to zero when no information is present in the data block. While evaluating the possibility of a specific BDS code, if any of the data blocks violate this rule, this BDS code is discarded.

In the structures of some Comm-B types, some bits can be reserved (or not used). These bits are required to be zeros at all times. If any of these predefined zero bits are set to one, the corresponding BDS code is also discarded.

The heuristic logic also checks parameter values that are decoded for different assumed BDS codes in a parallel fashion. These values need to be within their physical boundaries. For example, in BDS 60, the Mach number cannot be higher than one (for commercial aircraft). If a decoded Mach number is larger than one, BDS 60 is rejected as a BDS code. All evaluations are performed for each BDS code. Deterministic conditions for each parameter in all seven common BDS codes are listed in Table V.

2) Probabilistic Identification: Of all available surveillance messages, BDS 50 (track and turn report) and BDS 60 (heading and speed report) messages are most frequently interrogated. However, they have very similar structures, as shown in Table V. Due to this similarity, some messages can be considered as both BDS 50 and 60 after the heuristic logic. In order to differentiate these two codes, the probabilistic step is designed.

The principle is to construct a probability density function (PDF) using current aircraft states that are observed from ADS-B data. The probabilities of BDS 50 and 60 are computed with this function. The BDS code that results in a higher probability is considered to be the correct one.

From ADS-B, we use aircraft ground speed (v_g) and track angle (χ) to construct the *x* and *y* components of the velocity. They are treated as the means for the joint probability density function. These two components, denoted as v_{gx} and v_{gy} , as calculated as:

$$
v_{gx} = v_g \cdot \sin \chi
$$

\n
$$
v_{gy} = v_g \cdot \cos \chi
$$
 (5)

Then, we construct a bi-variate normal probability density function considering these values as mean values. To simplify the problem, we assume there is no correlation between the *x* and *y* speed components, the function of the probability (the Gaussian PDF without normalization) can be expressed as follows:

$$
p(v_{ax}, v_{ay}) = \exp\left\{-\frac{1}{2}\left[\frac{(v_{ax} - (v_{gx} - v_{wx}))^2}{\sigma_{ox}^2} + \frac{(v_{ay} - (v_{gy} - v_{wy}))^2}{\sigma_{oy}^2}\right]\right\}
$$
(6)

where v_{ax} and v_{ay} are the *x* and *y* components of airspeed. v_{wx} and v_{wy} are wind speed components. σ_{vx}^2 and σ_{vy}^2 are the variances. Accurate wind speed is not required for the purpose of BDS identification. For example, it can be obtained using weather forecast data. In case the wind information is not available or at calm wind conditions, these terms may be

BDS	Bits	Parameter	Logic rule	Value rule		
	$1 - 8$	BDS	Bits equal to 00010000			
10	$10 - 14$	Reserved	Bits must all be zeros			
17	7	BDS 20 enabled ^I	Bit equals to 1			
	$29 - 56$	Reserved	Bits must all be zeros			
20	$1 - 8$	BDS	Bits equal to 00100000			
	$9 - 56$	Call sign		Only contain 0-9, A-Z, or space		
	$1 - 8$	BDS	Bits equal to 00110000			
30	$29 - 30$	Threat type	Bits must not equal to 11			
	$16 - 22$	ACAS III		Less than 48		
	$1:2-13$ $^{\text{II}}$	MCP/FCU selected altitude	Status and value bits consistent III			
	$14:15-26$	FMS selected altitude	Status and value bits consistent			
40	$27:28-39$	Barometric pressure setting	Status and value bits consistent			
	$40 - 47$	Reserved	Bits must all be zeros			
	$52 - 53$	Reserved	Bits must all be zeros			
	$1:2-11$	Roll angle	Status and value bits consistent	Between -60 and 60 (degrees)		
	$12:13-23$	True track angle	Status and value bits consistent			
50	$24:25-34$	Ground speed	Status and value bits consistent	Between 0 and 600 (knots)		
	$35:36-45$	Track angle rate	Status and value bits consistent			
	$45:46-56$	True airspeed	Status and value bits consistent	Between 0 and 500 (knots)		
	$1: 2-12$	Magnetic heading	Status and value bits consistent			
	$13:14-23$	Indicated airspeed	Status and value bits consistent	Between 0 and 500 (knots)		
60	$24:25-34$	Mach number	Status and value bits consistent	Between 0 and 1		
	$35:36-45$	Barometric vertical rate	Status and value bits consistent	Between -6000 and 6000 (feet/minute)		
	$46:47-56$	Inertial vertical rate	Status and value bits consistent	Between -6000 and 6000 (feet/minute)		

TABLE V STATUS BIT AND PARAMETER BITS

^I BDS 20 is the code that has to be enabled for all transponders to provide the minimum Mode-S capabilities.

II This format $b1:b2-b3$ indicates the status bit at $b1$ with value bits from $b2$ to $b3$.

 III Consistency indicates that when status bit is zero, all value bits must also be zeros.

Fig. 7. Example of BDS 50/BDS 60 identification.

set to zero. However, this assumption may affect the accuracy of the identification.

For the variances, both are empirically set to be 20 knots in our experiments. This choice is based on the common magnitude of accuracy in speed which is included in ADS-B data [15]. Using Equation 6, the probabilities of BDS 50 and BDS 60 are compared. Finally, the corresponding BDS code to the higher $p(v_{ax}, v_{ay})$ value is accepted. Fig. 7 shows an example of the identification based on the possible speeds decoded as BDS 50 and BDS 60.

V. EXPERIMENT AND VALIDATION

In this section, we use three different datasets to verify, validate, and analyze the Mode-S Comm-B replies. The first

Fig. 8. Heuristic BDS detection, based on Comm-B replies received from 12:00 to 13:00 UTC, September 7, 2017, Delft, the Netherlands.

two datasets are from the same hour, with one collected by our receiver (located at the Delft University of Technology) and the other provided by the Air Traffic Control the Netherlands (LVNL). The last one is a global dataset provided by the ADS-B Exchange receiver network.

A. Experiment 1 - Examination the Heuristic Logic

This experiment focuses on examining the effectiveness of the HP process of the BDS identification presented in this paper. A one-hour dataset collected by our receiver is used as a test set (from 12:00 to 13:00 UTC, September 7, 2017). The result of the heuristic logic is shown in Fig. 8. In total, 1.5 million Comm-B replies are received during this one-hour period. The most common three BDS codes are 40, 50, and 60.

TABLE VI VALIDATION STATISTICS OF MODE-S BDS CODE IDENTIFICATION USING PROPOSED *HP Process*

Result	Heuristic logic	Probabilistic id
Correct	235,772 (93.169%)	247.192 (97.682%)
Unidentifiable	17,281 (6.829%)	5847 (2.311%)
Incorrect	$6(0.002\%)$	20(0.008%)

In each bar plot of Fig. 8, the hatch pattern marks the corrupted messages.

In this figure, we can see that around 20% of the messages are not identified (with unknown status). It is also noticeable that around 90% of these unidentified messages are corrupted.

For all other messages with a BDS code identified, BDS 10, 17, and 20 combined represent around 5% of the total number of messages. BDS 40 and 60 account for about 27% each, and BDS 50 around 16%. Due to the similar message structures, around 4% of the messages identify as both BDS 50 and 60 codes. Less than 0.1% of messages are BDS 30. Since BDS 30 encodes emergency status and ACAS resolutions, it is expected that this category is less frequently interrogated during normal operations.

B. Experiment 2 - Identification Accuracy

To examine the accuracy of the HP process for BDS identification, we obtain a true reference dataset of the Dutch airspace from the Air Traffic Control the Netherlands (LVNL). This data is collected in the same hour as the previous test dataset. The original data is presented in ASTERIX format, from which we extract the raw messages. As such, raw messages only contain messages that are of BDS 40, 50 and 60 (Enhanced Mode-S only). Raw messages for other BDS types are not included in this ASTERIX dataset.

After applying both heuristic and probabilistic steps, we compute the accuracy at each stage. The results are shown in Table VI. In total, 253,059 messages are used. Unidentifiable messages refer to the ones with multiple BDS codes.

Using only heuristic logic, we achieve a correctness rate of 93.2% for a total of 250 thousand messages. 6.8% of the messages have been identified with more than one BDS code. Using additional probabilistic identification, we can further reduce the rate of uncertain messages to 2.3%, and increase the success rate to 97.7%.

In each step, the number of incorrectly decoded messages is almost negligible. We have counted that only 6 (heuristic step) and 20 (probabilistic step) out of 253,059 messages were incorrectly identified. Regarding the significance of the probabilistic step, we found that among all 11,434 identifiable messages (with BDS 50 and BDS 60 combination) only 14 errors occurred. Upon further investigation, we notice that errors happen often when part of the flight information is missing from the message.

C. Experiment 3 - Analysis of the Global Data

Using multiple receivers provided by the ADS-B Exchange receiver network, we are able to examine the decoding performance of Mode-S Comm-B data across different regions

of the world. A set of data consisting of one hour (local time 10:00 AM to 11:00 AM, on 24 or 25 August of 2018) is collected at eight different locations around the world. Using the method proposed in this paper, we decode the messages and show the composition of Mode-S responses for these different regions. In Fig. 9, the percentages of BDS codes are illustrated.

Mode-S data transmitted to ADS-B Exchange network is filtered by default. Only messages with *s*² state in Fig. 5 are kept. These are Comm-B messages with ICAO addresses that appear in ADS-B data. In Table VII, the exact percentages of BDS codes in Comm-B responses are listed. In all eight regions, the unidentifiable messages are between 1% to 6% of received messages. This result is in line with the test performed using our own receiver after corrupted messages are discarded (as shown in Fig. 8). In most of these regions, common interrogated BDS codes are BDS40, BDS50, and BDS60, with the exception of Dallas in the USA and Tel Aviv in Israel. The diversity of interrogations is one of the most challenging elements in the inference process. However, the proposed methods in this paper are able to cope with interrogation variations and produce a large percentage of identifiable messages.

VI. DISCUSSIONS

This study aims at efficiently and accurately making use of open aircraft surveillance data. The focus has been on designing a path to decipher Comm-B replies that are part of the Mode-S enhanced surveillance and cannot be decoded directly. The experiments and validations show a high accuracy of the identification process and error detection strategy that are proposed in this paper. These refined identification methods provide the possibility to conduct aircraft performance studies with better accuracy when using open aircraft surveillance data.

Accurate identification of BDS 50 and 60 enables precise airspeed observations from Comm-B data. One possible extension of this research is its potential contribution to atmospheric modeling. By combining ADS-B ground speed with Mode-S true airspeed, wind and temperature can be computed. With globally around ten thousand aircraft airborne at any given time, a large number of meteorological measurements can be provided. This idea has been proposed in earlier research [16], where data are supplied by air traffic controllers. The methods and tools from this paper allow anyone to gain such ability. For example, in our earlier study [17], wind and air temperature were derived based on this open data, which demonstrates the use of aircraft for weather observations.

Fig. 3 and 6 in earlier sections show the percentage of corrupted messages found in this study. This strongly suggests a frequency congestion problem on the 1090 MHz channel. Of all the ADS-B and Comm-B messages that we received, more than half are corrupted. Based on the percentage of errors, severe message corruption during daily operations is shown. In Fig. 10, the percentage of corrupted ADS-B messages and the number of aircraft flying in a 24-hour period are plotted.

It can be seen that during the night time, when the airspace is less saturated, the percentage of corrupted

Fig. 9. Statistics of BDS codes in Comm-B replies at different geographic regions.

TABLE VII MODE-S COMM-B RESPONSE SUMMARY IN DIFFERENT REGIONS

Country	City	BDS ₁₀	BDS17	BDS ₂₀	BDS30	BDS40	BDS50	BDS60	EMPTY	MULTIPLE	UNKNOWN
Australia	Oueenland	0.3%	0.4%	1.1%	÷	31.7%	31.5%	32.9%	0.1%	0.3%	1.7%
Australia	Wollongong	0.3%	0.2%	0.5%	\sim	28.9%	31.6%	32.5%	0.3%	0.6%	5.0%
Belguim	Brussels	2.3%	$.2\%$	4.0%	0.0%	36.0%	17.1%	34.5%	0.7%	1.2%	3.0%
Israel	Tel Aviv	4.5%	.6%	28.5%	\sim	22.4%	21.4%	16.2%	0.7%	1.1%	3.7%
Russia	Pushkino	2.3%	3.1%	8.6%	0.8%	28.6%	27.2%	20.5%		4.2%	4.6%
UK	London	$.0\%$	1.7%	6.6%	0.0%	32.8%	20.6%	30.9%	0.5%	2.6%	3.1%
USA	Dallas	20.8%	59.2%	10.9%	-		$.0\%$	<u>т</u>	6.5%		1.6%
USA	San Francisco	9.3%	7.4%	6.5%	\sim	23.0%	19.9%	23.0%	4.1%	1.3%	5.5%

Fig. 10. Percentage of corrupted ADS-B messages and number of aircraft per hour on May 30, 2018, in Delft, the Netherlands.

messages decreases. However, the decrease in corrupted messages (∼10%) is not proportional to the decrease in the number of flights (∼65%). These corrupt messages not only present a challenge for obtaining more accurate data but also indicate a constant frequency congestion in busy airspace.

Other than third-party researchers, air traffic controllers may also benefit from the methods proposed in this paper. With a high level of frequency congestion as shown in Fig. 8, air traffic controllers tend not to over-interrogate. This creates a dilemma where an ATC has to reduce the number and frequency of interrogations, but at the same time, more information is preferable for making better traffic control decisions. Using the proposed BDS identification process, one ATC can intercept replies that originate from interrogations by another

ATC center. As one example, based on the one-hour reference dataset, we found that the replies from one specific secondary radar from the air traffic control of the Netherlands account for only around 17% of the total number of interrogation replies. The proposed identification method can therefore significantly increase the amount of information available to each ATC directly, even for aircraft that are outside of their airspace.

Finally, it is worth emphasizing that Mode-S Comm-B messages are passive replies originating from SSR interrogations. For oceanic and remote areas where secondary surveillance radars are not present, only ADS-B data is available. In this paper, not all BDS codes are addressed. Instead, only the types related to ELS and EHS are investigated. However, this is the situation for many regions around the world, as illustrated in Fig. 9 using the global Mode-S data.

VII. CONCLUSIONS

In this paper, we propose a set of inference methods that allow third-party observers to identify and decode Mode-S Comm-B replies. The inference is based solely on surveillance replies without any knowledge of Mode-S interrogation. This paper contributes to a missing area of knowledge on handling interrogation-based surveillance data. It gives researchers broader access to accurate aircraft state updates that are transmitted through Enhanced Mode-S. The implementation is based on existing low-cost commercial off-the-shelf ADS-B receivers, with no additional hardware. Using a reference dataset, the proposed process reveals a correctness rate of around 97.7%, with 2.3% unidentifiable and 0.008% error for Enhanced Mode-S messages.

Furthermore, the process proposed in this paper is also likely to be beneficial for air traffic controllers, since it enables the ability to collect more data without the need for increasing interrogation frequency. Finally, we have made the decoding process into an open-source programming library, *pyModeS*, which includes the decoding and inferences discussed in this paper. We hope this will also enable other researchers to make use of this valuable aircraft surveillance data source.

ACKNOWLEDGMENTS

We would like to thank Air Traffic Control the Netherlands (LVNL) for providing us the Enhanced Mode-S validation dataset. We would also like to thank James Stanford from ADS-B Exchange for providing the data for global data analysis.

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